PCC Vivace: Online-Learning Congestion Control

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Internet Congestion Control



Internet Congestion Control



Packet loss / RTT increment indicates congestion

Internet Congestion Control



Strong Assumptions Cause Problem



Abstract assumption cannot capture Internet complexity

PCC Utility Framework



PCC Rate Control



Sender selfishly maximizes its own utility (*online learning* in *non-cooperative* game)

PCC Design Challenges



Requirements for consistently high performance:

- Capture application performance objectives
- Guarantee equilibrium with multiple competing senders

- Guarantee reaching equilibrium upon convergence
- Rapidly adapt to network dynamics

PCC Allegro

Heuristic Loss-based utility function rate control $u_i(x_i) = T_i \cdot Sigmoid_{\alpha}(L_i - 0.05) - x_i \cdot L_i$ *Fixed rate* $T_i = x_i(1 - L_i)$ change step size Sigmoid_{α}(y) = $\frac{1}{1+e^{\alpha y}}$ No latency-awareness Slow convergence Slow reaction to network changes Can cause bufferbloat

[Dong et al. NSDI 2015]



PCC Vivace

• Leveraging powerful tools from online learning theory

New utility function framework

- Latency-awareness
- Strictly concave \Rightarrow Equilibrium guarantee
- Flexibility among senders

New control algorithm

- Gradient-ascent \Rightarrow Convergence speed/stability
- Deals with measurement noise

PCC Vivace





$$\frac{u(x+\delta) - u(x-\delta)}{2\delta} \to \alpha \cdot \Delta x$$

Techniques to deal with measurement noise:

- Linear regression
- RTT gradient low-pass filter
- Double check







"No-regret" guarantee: A Powerful lens for analysis

Evaluation

- Implementation
 - UDT-based user-space implementation
 - Emulab experiments, Amazon EC2 experiments
 - User-space PCC proxy for video streaming
- Protocols for comparison
 - TCP variants (TCP CUBIC, TCP Illinois, TCP Vegas, *etc.*)
 - BBR
 - PCC Allegro
 - PCC Vivace

Vivace Utility Performance

• Latency awareness (100Mbps, 30ms RTT Emulab bottleneck link)



< 2ms inflation in all cases 90% smaller than BBR under 2BDP

PCC Vivace: Online-Learning Congestion Control

Vivace Rate Control Performance

 Rapid reaction to network changes (10–100Mbps, 10–100ms RTT, 0–1% random loss)
 PCC Allegro PCC Vivace



Vivace upon Convergence

• Fair equilibrium (100Mbps, 30ms RTT, 75KB buffer)



TCP Friendliness



Insights from Learning-Theoretic Tools

• Flexible equilibrium state with heterogeneous senders

$$u(x_i, L_i) = x_i - c_i x_i \left(\frac{1}{1 - L_i} - 1\right)$$
 $c_i = \frac{C}{x_i^*}$



Limitation in Extremely Dynamic Networks

• LTE (Mahimahi emulator [Netravali et al. ATC 2015])



https://www.youtube.com/watch?v=Y3IzuCdwdUo&t=27s
(Demo comparing PCC with UDP and TCP video streaming)

PCC In Action

- Open source release on GitHub (<u>https://github.com/pccproject</u>)
 - UDP implementation used in experiments presented here
 - QUIC implementation with **Google**
 - Pantheon implementation for test purpose
 - Kernel implementation in the works
- VACC variant of PCC by 处 HUAWEI and 🔘 vodafone
 - Kernel implementation with optimizations for video over LTE
 - Ongoing research project with successful field tests

Conclusion

- PCC Vivace: Leveraging no-regret learning for congestion control
 - Consistent high performance as PCC Allegro
 - Latency awareness, mitigated bufferbloat (latency inflation, congestion loss)
 - Provably fair, yet also flexible equilibrium convergence
 - Fast and stable convergence, even with changing network conditions
 - Improved TCP friendliness, safer to deploy
- Thanks for generous project support



Thanks!





Vivace Utility Performance

• Latency awareness (100Mbps, 30ms RTT Emulab bottleneck link)



< 0.5% loss with 13.5KB buffer 55% smaller than TCP CUBIC

Vivace Rate Control Performance

 Convergence Speed/Stability Tradeoff (100Mbps, 30ms RTT, 75KB buffer)



Insights from No-Regret Guarantee

Random loss tolerance vs. Congestion loss (8Mbps, 25KB per-flow share)



Performance in Real-World



Limitation of No-Regret

"Sender's choices of rates are asymptotically (across time) no worse, utility-wise, than sending at what would have been (in hindsight) the best *fixed* rate"

Still make sense in highly dynamic environment?